MPhil Advanced Computer Science
Topics in Logic and Complexity

Lent 2022 — Anuj Dawar

Exercise Sheet 2

1. In the lecture we saw an illustration of a construction to show that *acyclicity* of graphs is not definable in first-order logic. Write out a proof of this result.

   Prove that *acyclicity* is not definable in $\mathsf{Mon}.\Sigma^1_1$. Is it definable in $\mathsf{Mon}.\Pi^1_1$?

2. Prove (using Hanf's theorem or otherwise) that 3-colourbility of graphs is not definable in first-order logic.

   Graph 3-colourability (and, indeed, 2-colourability) are definable in $\mathsf{Mon}.\Sigma^1_1$. Can you show they are not definable in $\mathsf{Mon}.\Pi^1_1$? Are they definable in *universal second-order logic*?

3. Prove the lemma stated in the lecture that any formula that is positive in the relation symbol R defines a monotone operator.

4. Prove that the formula of LFP $\neg[\mathbf{lfp}_{R,\mathbf{x}}\neg\phi(R/\neg R)](\mathbf{x})$, where $\phi(R/\neg R)$ denotes the result of replacing all occurrences of R in ϕ by $\neg R$, defines the greatest fixed point of the operator defined by ϕ.

5. In the lectures, we saw how definitions by simultaneous induction can be replaced by a single application of the **lfp** operator. In this exercise, you are asked to show the same for *nested* applications of the **lfp** operator.

   Suppose $\phi(\mathbf{x}, \mathbf{y}, S, T)$ is a formula in which the relational variables S (of arity s) and T (of arity t) only appear positively, and $\mathbf{x}$ and $\mathbf{y}$ are tuples of variables of length s and t respectively. Show that (for any t-tuple of terms $\mathbf{t}$) the predicate expression

   $$[\mathbf{lfp}_{S,\mathbf{x}}([\mathbf{lfp}_{T,\mathbf{y}}\phi](\mathbf{t}))]$$

   is equivalent to an expression with just one application of **lfp**.

6. Consider a vocabulary consisting of two unary relations P and O, one binary relation E and two constants s and t. We say that a structure $\mathbb{A} = (A, P, O, E, s, t)$ in this vocabulary is an *arena* if $P \cup O = A$ and $P \cap O = \emptyset$. That is, P and O partition the universe into two disjoint sets.

   An arena defines the following game played between a *player* and an *opponent*. The game involves a *token* that is initially placed on the element s. At each move, if the token is currently on an element of P it is *player* who plays and if it is on an element of O, it is *opponent* who plays. At each move, if the token is on an element a, the one who plays choses an element b such that $(a, b) \in E$ and moves the token from a to b. If the token reaches t at any point then *player* has won the game.

We define GAME to be the class of arenas for which *player* has a strategy for winning the game. Note that in an arena $\mathbb{A} = (A, P, O, E, s, t)$, *player* has a strategy to win from an element a if *either* $a \in P$ and there is some move from a so that *player* still has a strategy to win after that move *or* $a \in O$ and for every move from a, *player* can win after that move.

(a) Give a sentence of LFP that defines the class of structures GAME.

We say that a collection $\mathcal{C}$ of decision problems is *closed under logarithmic space reductions* if whenever $A \in \mathcal{C}$ and $B \leq_L A$ (i.e. B is reducible to A by a logarithmic-space reduction) then $B \in \mathcal{C}$.

The class of structures GAME defined above is known to be P-complete under logarithmic-space reductions.

(b) Explain why this, together with (a) implies that the class of problems definable in LFP is *not* closed under logarithmic-space reductions.

7. Give a sentence of LFP that defines the class of linear orders with an even number of elements.

8. The *directed graph reachability problem* is the problem of deciding, given a structure (V, E, s, t) where E is an arbitrary binary relation on V, and $s, t \in V$, whether (s, t) is in the reflexive-transitive closure of E. This problem is known to be decidable in NL.

   Transitive closure logic is the extension of first-order logic with an operator **tc** which allows us to form formulae

$$\phi \equiv [\mathbf{tc}_{\mathbf{x},\mathbf{y}}\psi](\mathbf{t}_1, \mathbf{t}_2)$$

   where $\mathbf{x}$ and $\mathbf{y}$ are k-tuples of variables and $\mathbf{t}_1$ and $\mathbf{t}_2$ are k-tuples of terms, for some k; and all occurrences of variables $\mathbf{x}$ and $\mathbf{y}$ in ψ are bound in ϕ. The semantics is given by saying, if $\mathbf{a}$ is an interpretation for the free variables of ϕ, then $\mathcal{A} \models \phi[\mathbf{a}]$ just in case $(\mathbf{t}_1^{\mathbf{a}}, \mathbf{t}_2^{\mathbf{a}})$ is in the reflexive-transitive closure of the binary relation defined by $\psi(\mathbf{x}, \mathbf{y})$ on A^k.

   (a) Show that any class of structures definable by a sentence ϕ, as above, where ψ is first-order, is decidable in NL.

   (b) Show that, if K is an isomorphism-closed class of structures in a relational signature including $<$, such that each structure in K interprets $<$ as a linear order and

$$\{[\mathcal{A}]_< \mid \mathcal{A} \in K\}$$

   is decidable in NL, then there is a sentence of transitive-closure logic that defines K.

9. For a binary relation E on a set A, define its *deterministic transitive closure* to be the set of pairs (a, b) for which there are $c_1, \ldots, c_n \in A$ such

that $a = c_1$, $b = c_n$ and for each $i < n$, c_{i+1} is the *unique* element of A with $(c_i, c_{i+1}) \in E$.

Let DTC denote the logic formed by extending first-order logic with an operator **dtc** with syntax analogous to **tc** above, where $[\mathbf{dtc}_{\mathbf{x},\mathbf{y}}\psi]$ defines the deterministic transitive closure of $\psi(\mathbf{x}, \mathbf{y})$.

(a) Show that every sentence of DTC defines a class of structures decidable in L.

(b) Show that, if K is an isomorphism-closed class of structures in a relational signature including $<$, such that each structure in K interprets $<$ as a linear order and

$$\{[\mathcal{A}]_< \mid \mathcal{A} \in K\}$$

is decidable in L, then there is a sentence of DTC that defines K.

10. The *structure homomorphism problem* for a relational structure σ (with no function or constant symbols) is the problem of deciding, given two σ-structures $\mathbb{A}$ and $\mathbb{B}$ whether there is a homomorphism $\mathbb{A} \longrightarrow \mathbb{B}$.

    (a) Show that if σ contains only unary relations, then the structure homomorphism problem for σ is decidable in polynomial time.

    (b) Show that if σ contains a relation of arity 2 or more, then the the structure homomorphism problem for σ is NP-complete.

11. Schaefer's theorem (Handout 5, page 12) gives six conditions under which CSP($\mathbb{B}$) is in P, for $\mathbb{B}$ a structure on domain $\{0, 1\}$. For each of the six conditions, show that indeed any CSP($\mathbb{B}$) is in P.

    For the first five conditions, it is also the case that CSP($\mathbb{B}$) is definable in LFP. Prove this.

12. We saw (Handout 6, page 3) that for any $\mathbb{B}$, CSP($\mathbb{B}$) is definable in MSO. Write out an MSO formula for 3-SAT as defined on page 9 of Handout 5.

13. Prove that if CSP($\mathbb{B}$) has bounded width, it is definable in LFP.

14. We saw in the lecture that CSP(K_2) has width 3. Prove that CSP(K_3) does not have width 3. (*Hint*: Consider the graph K_4).