

8 Hoare Logic and Model Checking (jp622)

Consider commands  $C$  composed from assignments  $X := E$  (where  $X$  is a program variable, and  $E$  is an arithmetic expression), heap allocation  $X := \text{alloc}(E_1, \dots, E_n)$ , heap assignment  $[E_1] := E_2$ , heap dereference  $X := [E]$ , disposal of heap locations  $\text{dispose}(E)$ , the no-op **skip**, sequencing  $C_1; C_2$ , conditionals **if**  $B$  **then**  $C_1$  **else**  $C_2$  (where  $B$  is a boolean expression), and loops **while**  $B$  **do**  $C$ . **null** is 0.

Recall the separation logic partial list representation predicates:

$$\begin{aligned} \text{plist}(t, [], u) &= (t = u) \wedge \text{emp} \\ \text{plist}(t, h :: \alpha, u) &= \exists y. ((t \mapsto h) * ((t + 1) \mapsto y) * \text{plist}(y, \alpha, u)) \end{aligned}$$

Circular lists can be represented by  $\text{clist}(t, \alpha) = \text{plist}(t, \alpha, t) \wedge (\alpha = [] \Rightarrow t = \text{null})$ .

(a) Assuming  $\vdash \{P_1\} C_1 \{Q_1\}$  and  $\vdash \{P_2\} C_2 \{Q_2\}$ :

(i) explain precisely why  $\vdash \{P_1 * P_2\} C_1; C_2 \{Q_1 * Q_2\}$  [2 marks]

(ii) give a counterexample to  $\vdash \{P_1 \wedge P_2\} C_1; C_2 \{Q_1 \wedge Q_2\}$ . [1 mark]

(b) Give a proof outline for the following circular list ‘next’ triple:

$$\{\text{clist}(X, t :: \alpha)\} X := [X + 1] \{\text{clist}(X, \alpha ++ [t])\} \quad [3 \text{ marks}]$$

(c) Give a loop invariant (no need for a proof outline) for the following circular list ‘length’ triple:]

$$\begin{aligned} &\{\text{clist}(X, \alpha)\} \\ &\text{if } X = \text{null} \text{ then } Y := 0 \\ &\text{else } (Z := [X + 1]; Y := 1; \text{while } Z \neq X \text{ do } (Z := [Z + 1]; Y := Y + 1)) \\ &\{\text{clist}(X, \alpha) * Y = \text{length}(\alpha)\} \end{aligned}$$

[3 marks]

(d) Give a loop invariant (no need for a proof outline) for the following triple for a ‘previous’ operation on non-empty circular lists:

$$\begin{aligned} &\{\text{clist}(X, \alpha ++ [t])\} \\ &Z := X; Y := [X + 1]; (\text{while } Y \neq X \text{ do } (Z := Y; Y := [Y + 1])); X := Z \\ &\{\text{clist}(X, t :: \alpha)\} \end{aligned}$$

[4 marks]

(e) Give a loop invariant (no need for a proof outline) for the following triple for a ‘dial to minimum’ operation on non-empty circular lists:

$$\begin{aligned} &\{\text{clist}(X, \alpha_1 ++ [t] ++ \alpha_2) \wedge \text{sorted}(t :: \text{merge}(\text{sort}(\alpha_1), \text{sort}(\alpha_2)))\} \\ &Z := X; M := [X]; Y := [X + 1]; \\ &(\text{while } Y \neq Z \text{ do} \\ &\quad (N := [Y]; (\text{if } N < M \text{ then } X := Y \text{ else skip}); Y := [Y + 1])); \\ &\{\text{clist}(X, [t] ++ \alpha_2 ++ \text{reverse}(\alpha_1))\} \end{aligned}$$

[5 marks]

(f) Describe precisely *all* pairs of a stack and a heap that satisfy

$$\exists y, z. ((X \mapsto y * y \mapsto z * z \mapsto X) \wedge Y = 0)$$

[2 marks]