

Getting Started With Isabelle

Lecture I: Tour

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What is Isabelle?

a generic proof assistant ...

... based on a logical framework

a tool for mechanizing formalisms

a proof environment for its built-in logics:

- ZF set theory
- HOL (higher-order logic)
- many others including TLA & UNITY



Generic Features — Available to Many Logics

Theories declaring types, constants, etc. & inheriting from other theories

Flexible syntax including rewrite rules on abstract syntax trees

Order-sorted polymorphism to generalize results over related structures

Simplifier accepting conditional, permutative rewrite rules, ...

Classical reasoner to search for proofs using analytic rules

Features Specific to Isabelle/HOL

Proof libraries for integers, reals, lists, sets, cardinality, ...

Worked examples in semantics, security, concurrency, non-standard analysis, ...

Datatype definitions to model recursive types in functional programs

Inductive and Co-Inductive definitions to formalize semantics

Recursive functions defined over arbitrary well-founded relations

AND a link-up to SVC, the Stanford Validity Checker

Logical Reasoning in Isabelle/HOL

```
Goal "(ALL x. honest(x) & industrious(x) --> healthy(x)) & \
\ ~ (EX x. grocer(x) & healthy(x)) & \
\ (ALL x. industrious(x) & grocer(x) --> honest(x)) & \
\ (ALL x. cyclist(x) --> industrious(x)) & \
\ (ALL x. ~ healthy(x) & cyclist(x) --> ~ honest(x)) \
\ --> (ALL x. grocer(x) --> ~ cyclist(x))";
```

The command `Goal` states the formula to be proved.

```
by (Blast_tac 1);
```

The command `by` applies a `tactic` to the subgoals.

- Proved in zero seconds!
- `Blast_tac` is a powerful, generic tableau prover.

Set-Theoretic Reasoning in Isabelle/HOL

```
Goal "(INT i:I. A(i) Int B(i)) = \
\      (INT i:I. A(i)) Int (INT i:I. B(i))";
```

Here the goal is

$$\left(\bigcap_{i \in I} A_i \cap B_i \right) = \left(\bigcap_{i \in I} A_i \right) \cap \left(\bigcap_{i \in I} B_i \right)$$

```
by (Blast_tac 1);
```

- Blast_tac's default rules cover set theory and much more!
- You can insert new default rules.

Arithmetic Reasoning I: The Theory File

```
NatSum = Main +  
  
consts sum      :: "[nat=>nat, nat] => nat"  
  
primrec  
  "sum f 0          = 0"  
  "sum f (Suc n) = f(n) + sum f n"  
  
end
```

Theory `NatSum` extends `Main`, the standard parent.

Constant `sum` is declared with a curried function type.

Recursion equations make it a summation functional.

Arithmetic Reasoning II: Proving

$$1 + 3 + \cdots + (2n - 1) = n^2$$

```
Goal "sum (%i. Suc(i+i)) n = n*n";  
by (induct_tac "n" 1);
```

```
Level 1 (2 subgoals)  
sum (%i. Suc (i + i)) n = n * n  
  1. sum (%i. Suc (i + i)) 0 = 0 * 0  
  2. !!n. sum (%i. Suc (i + i)) n = n * n  
        ==> sum (%i. Suc (i + i)) (Suc n) = Suc n * Suc n
```

```
by Auto_tac;  
qed "sum_of_odds";
```

The tactic `induct_tac` applies structural induction, while `Auto_tac` simplifies and breaks up all subgoals.

Datatypes I: Specifying Boolean Expressions

```
datatype boolex = Const bool
                  | Neg boolex
                  | And boolex boolex

consts value :: "boolex => bool"
primrec
  "value (Const b) = b"
  "value (Neg b)   = (~ value b)"
  "value (And b c) = (value b & value c)"
```

Type `boolex` has three constructors.

Constant `value` maps these Boolean expressions to truth values.

It is declared *primitive recursive*.

Datatypes II: Specifying If-Expressions

```
datatype ifex = CIF bool | IF ifex ifex ifex

consts valif   :: "ifex => bool"
primrec
  "valif(CIF b)      = b"
  "valif(IF b t e) = (if valif b then valif t else valif e)"

consts bool2if :: "boolex => ifex"
primrec
  "bool2if(Const b) = CIF b"
  "bool2if(Neg b)   = IF (bool2if b) (CIF False) (CIF True)"
  "bool2if(And b c) = IF (bool2if b) (bool2if c) (CIF False)"
```

Functions `valif` and `bool2if` relate types `ifex`, `bool` and `boolex`.

Datatypes III: Proving bool2if Correct

```
Goal "valif (bool2if b) = value b";
by (induct_tac "b" 1);
```

Level 1 (3 subgoals)

```
valif (bool2if b) = value b
  1. !!bool. valif (bool2if (Const bool)) = value (Const bool)
  2. !!boolex.
    valif (bool2if boolex) = value boolex
    ==> valif (bool2if (Neg boolex)) = value (Neg boolex)
  3. !!boolex1 boolex2.
    [| valif (bool2if boolex1) = value boolex1;
       valif (bool2if boolex2) = value boolex2 |]
    ==> valif (bool2if (And boolex1 boolex2)) =
        value (And boolex1 boolex2)
```

```
by Auto_tac;
```

General Recursion I: Declaring QuickSort

```
Qsort = Sorting +
consts quickSort :: "('a::linorder) list => 'a list"

recdef quickSort "measure size"
simpset
  "simpset() addsimps [length_filter RS le_less_trans]"

"quickSort []      = []"
"quickSort (x#l) = quickSort [y:l. ~ x<=y] @
                    (x # quickSort [y:l. x<=y])"
```

Parent theory `Sorting` defines sorted and multiset.
Function `quickSort` is recursive in the size of its argument.

General Recursion II: A QuickSort Proof

```
Goal "multiset (quickSort xs) z = multiset xs z";
by (res_inst_tac [ ("u", "xs") ] quickSort.induct 1);

multiset (quickSort xs) z = multiset xs z
1. multiset (quickSort []) z = multiset [] z
2. !!x l.
   [| multiset (quickSort (filter (op <= x) l)) z =
      multiset (filter (op <= x) l) z;
      multiset (quickSort [y:l . ~ x <= y]) z =
      multiset [y:l . ~ x <= y] z |]
==> multiset (quickSort (x # l)) z =
      multiset (x # l) z

by Auto_tac;
```

Summary

Commands for managing an interactive proof:

- `Goal`: start it
- `by`: apply a tactic
- `qed`: name & store the proved theorem

Tactics for the reasoning itself

- `induct_tac i`: structural induction on subgoal *i*
- `Blast_tac i`: classical reasoning on subgoal *i*
- `Auto_tac`: tackle **all** subgoals

Theory file elements `consts`, `datatype`, `primrec`, ...